

Decentralized provision of active power by means of dynamic virtual power plants

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Research Cluster Smart Nord

Decentralised, agent based methods for control of medium and low voltage (smart) grids

Stability of medium and low voltage grids considering decentralized supply

Partners:

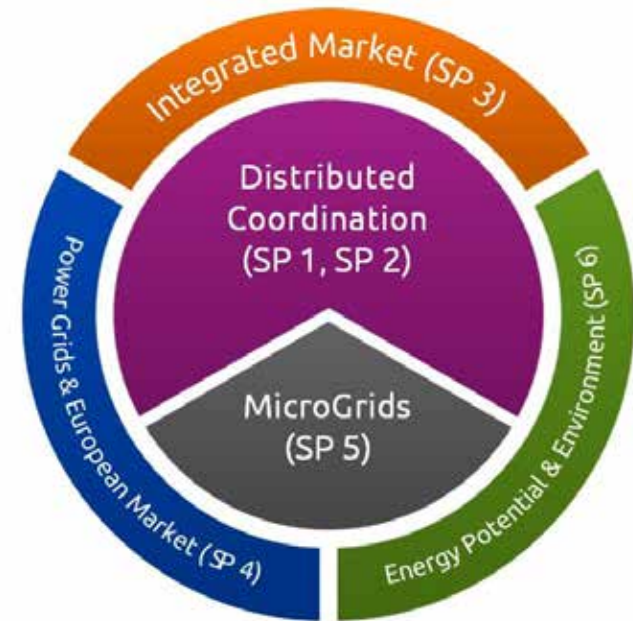
Uni Oldenburg, Uni Hannover,
TU Braunschweig, TU Clausthal,
NEXT ENERGY, OFFIS

Funding

4.1 M€ by Federal State of Lower Saxony

Duration

3/2012 – 2/2015



Dynamic Virtual Power Plants

DVPP

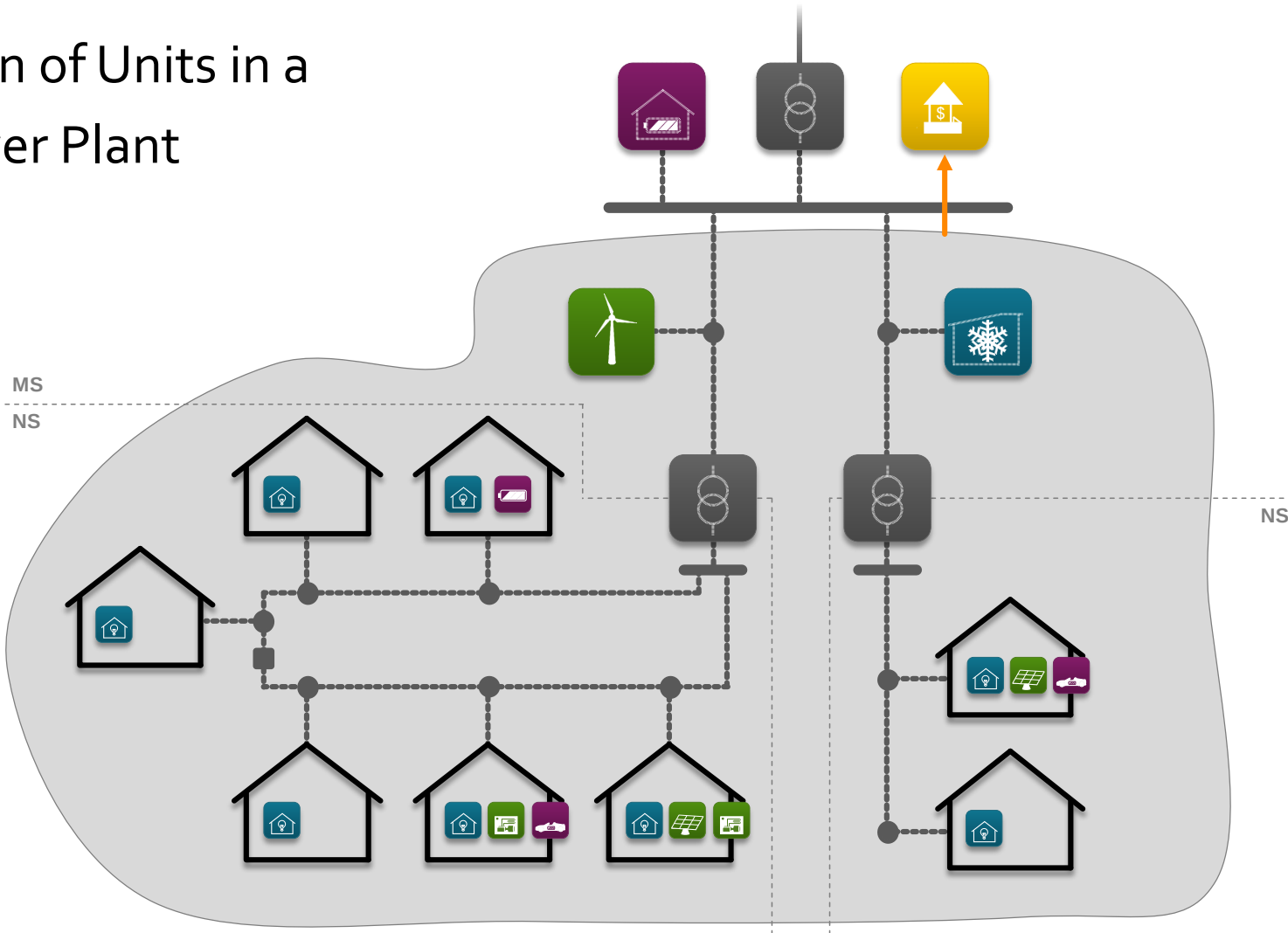
MOTIVATION AND CONCEPT

► Post EEG challenges for renewables

- ▶ **Renewable energy has to be traded at a power market**
- ▶ Requirement: Enhance supply of single renewable energy sources
 - ▶ Market barriers at the EEX (currently min. 100kW)
 - ▶ Reliability of supply regarding an agreed schedule (product)
- ▶ Smart Nord: *Dynamic, distributed* aggregation and control of units to adapt to
 - ▶ Demand traded at the market
 - ▶ Current set of active units in the system
 - ▶ Forecast of supply of single units

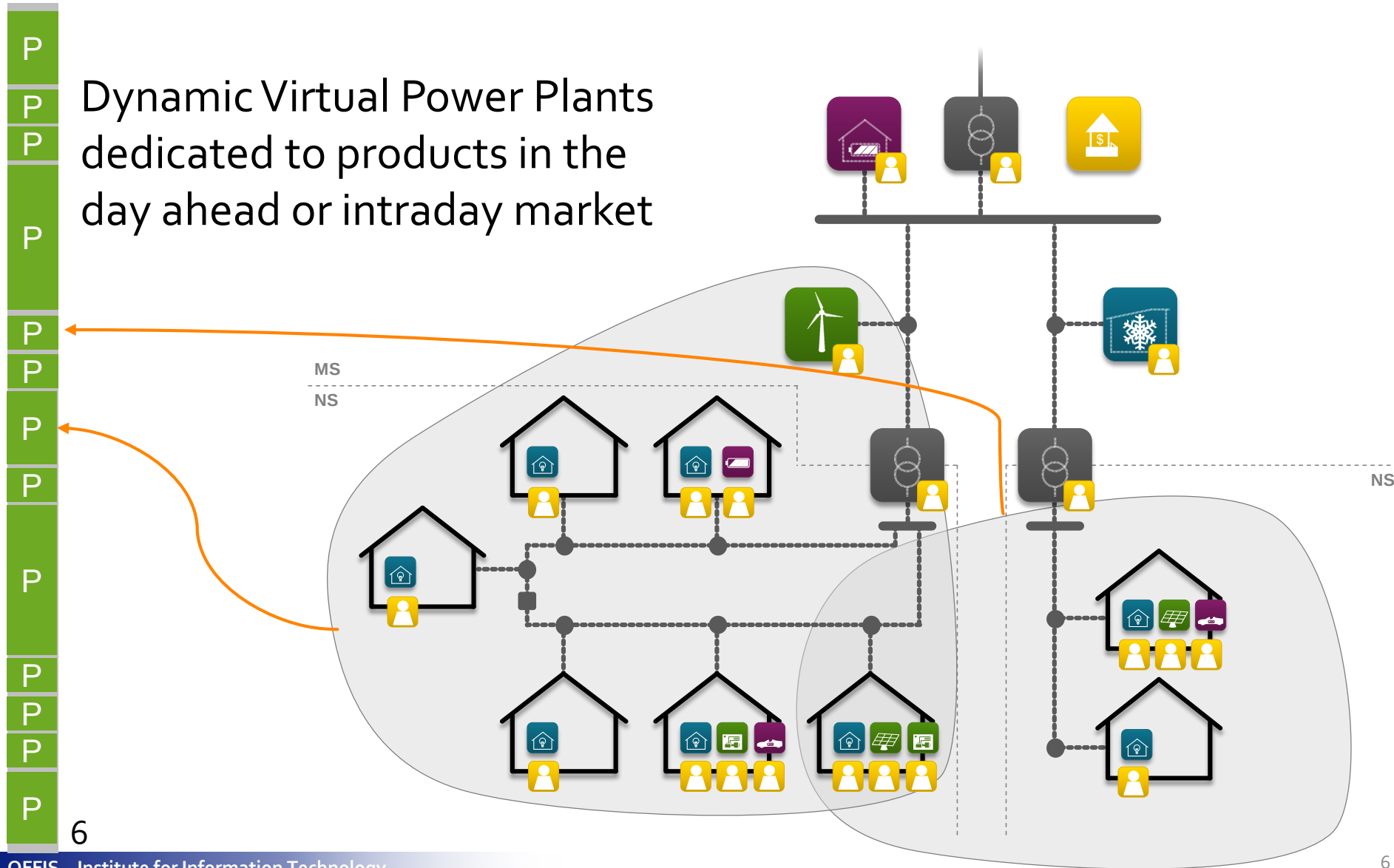
Virtual Power Plants (VPP)

Aggregation of Units in a Virtual Power Plant



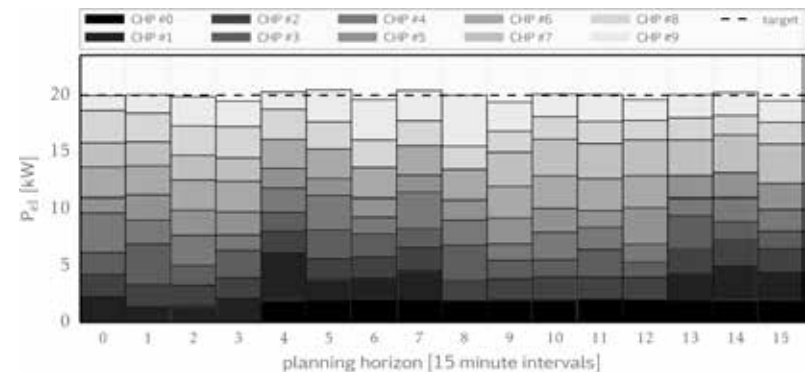
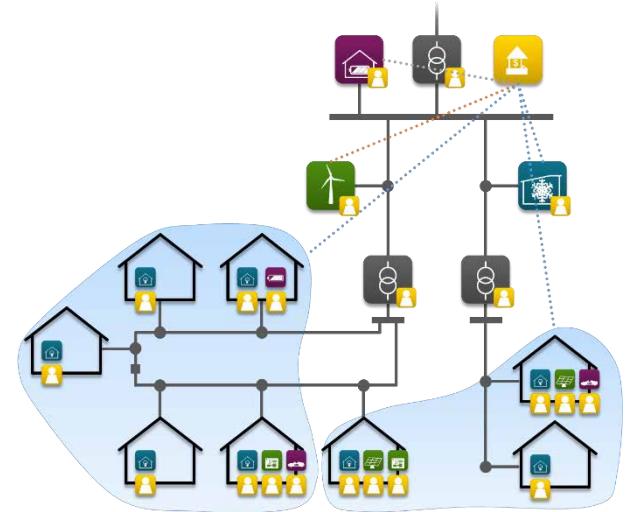
Dynamic Virtual Power Plants (DVPP)

Dynamic Virtual Power Plants
dedicated to products in the
day ahead or intraday market



Coordination of DVPPs

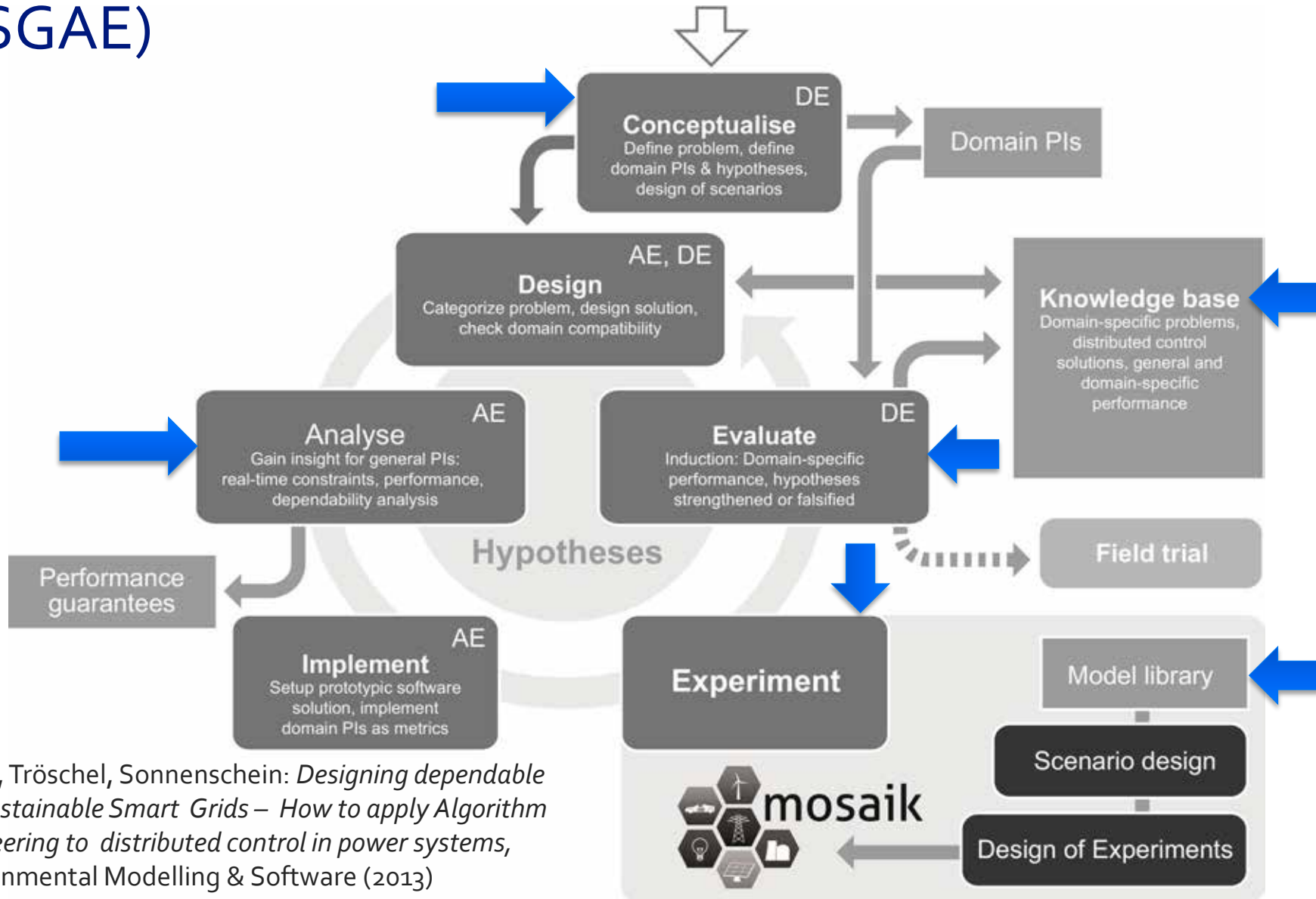
- ▶ DVPP setup:
Coalition formation of plants
for trading a specific power
product
- ▶ Predictive scheduling:
Determine a schedule which
 - ▶ guarantees the successful bid
 - ▶ and respects all individually
defined degrees of freedom
- ▶ Continuous scheduling:
Reschedule if unforeseen
changes occur
- ▶ Payoff division:
Distribute the gained revenues
just and reasonable to the
coalition members



▶ Decentralized, self-organized control

- ▶ Dynamic aggregation reacting to market conditions
 - ▶ Units possibly distributed to several parts of the grid:
no control centre, no micro grid approach
 - ▶ No knowledge about number of units:
Scalability of control methods
 - ▶ Ability to reconfigure dynamically in case of failures:
Robustness regarding failure of single units
- Decentralized, self-organized control
-
- ▶ Security

Smart Grid Algorithm Engineering (SGAE)



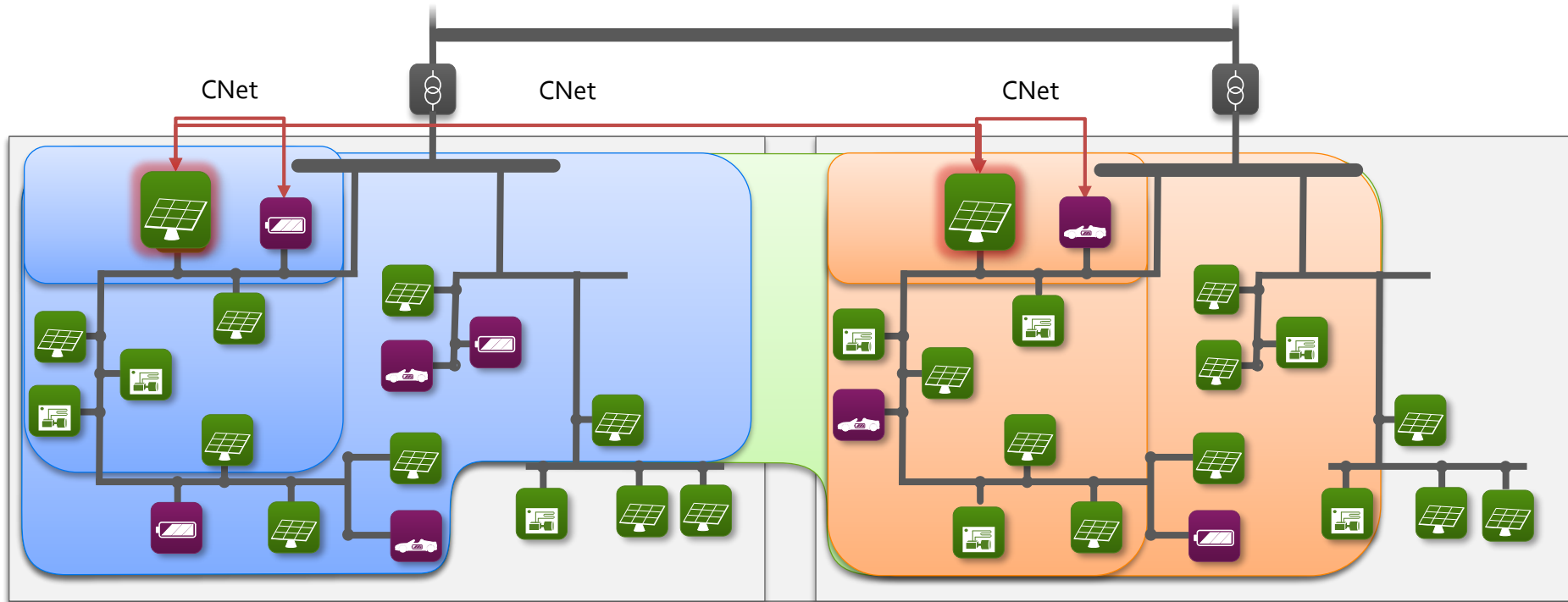
Nieße, Tröschel, Sonnenschein: *Designing dependable and sustainable Smart Grids – How to apply Algorithm Engineering to distributed control in power systems*, Environmental Modelling & Software (2013)

Forming coalitions to trade power at the market

DVPP SETUP

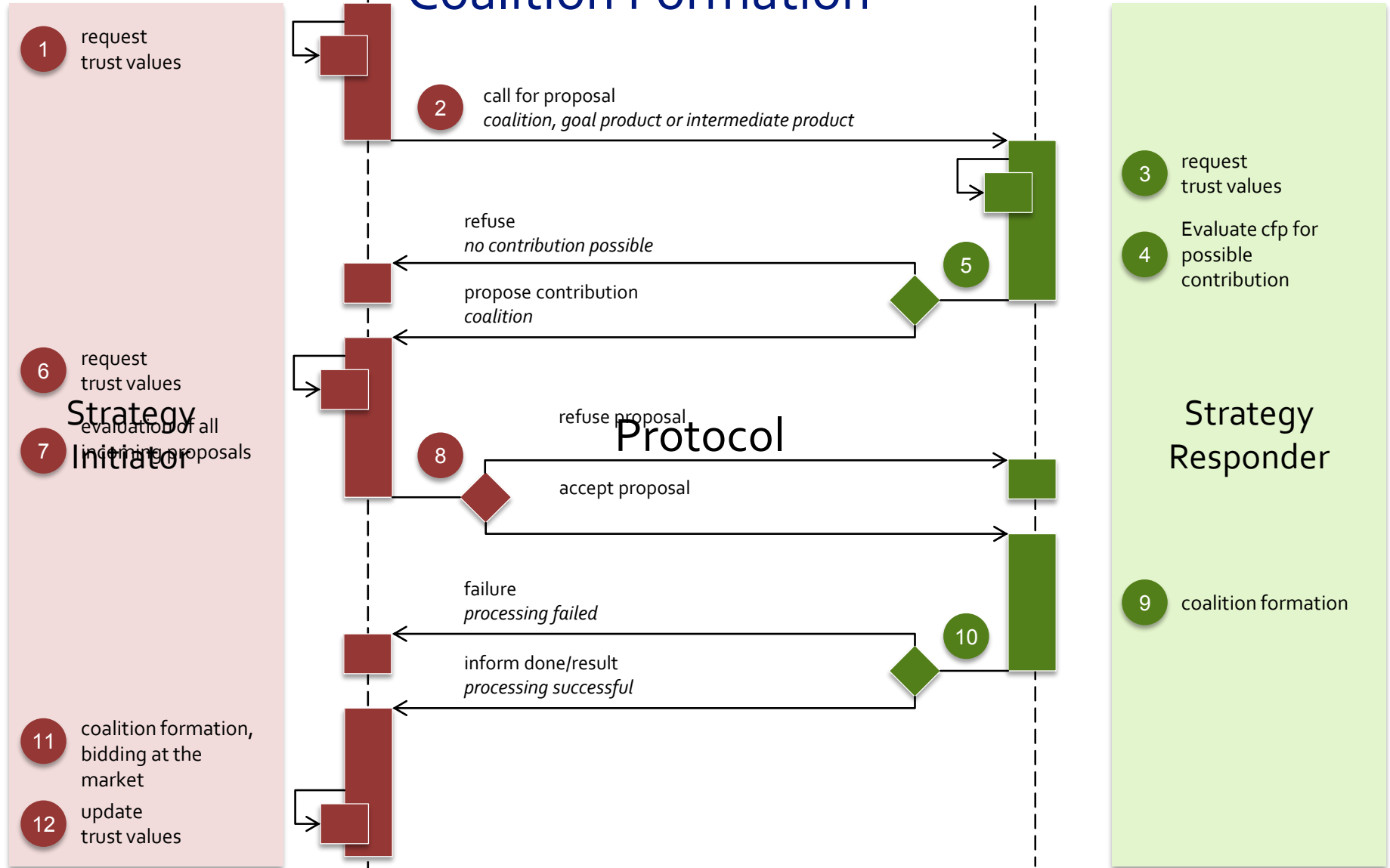
DVPP setup

Goal: deliver 100kW from 1 p.m. to 2 p.m.



- 1 Selection of Neighborhood
- 2 Coalition Formation
- 3 Expansion of Neighborhood

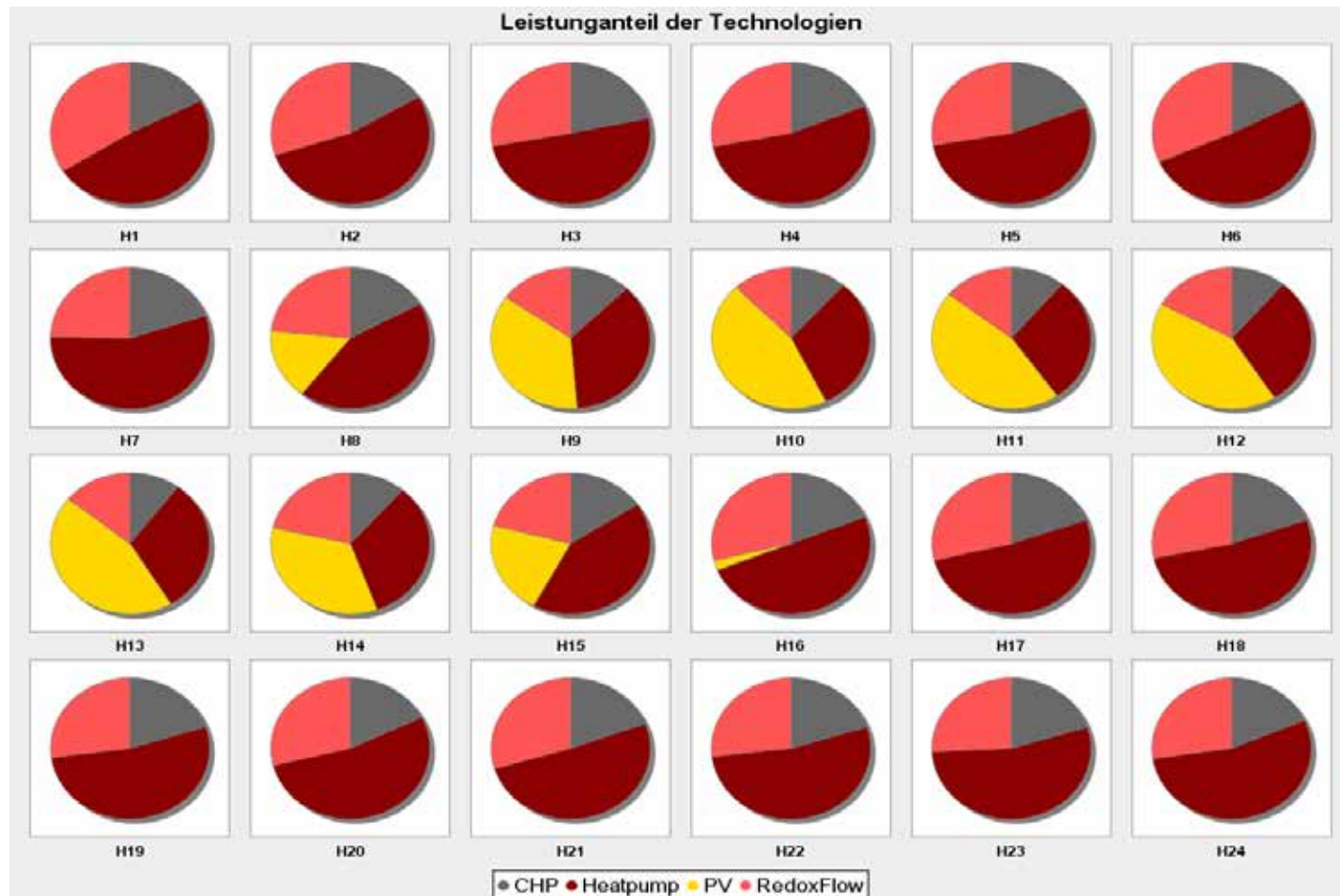
Coalition Formation



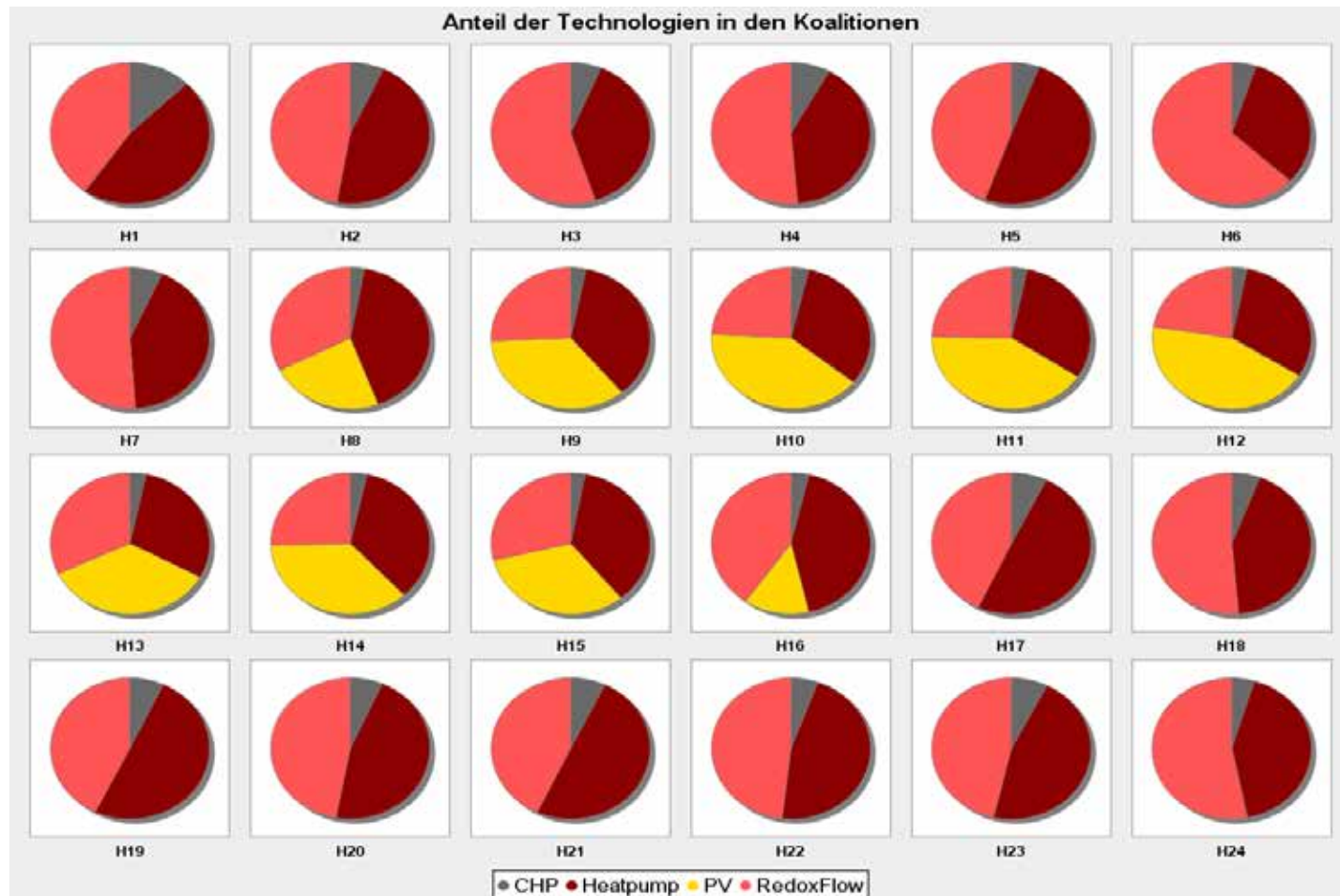
Evaluation Scenario 2030

- ▶ Simulated Units
 - ▶ 8806 households
 - ▶ 789 heat pumps
 - ▶ 1048 photovoltaic systems
 - ▶ 28 wind power generators
 - ▶ 122 CHPs
 - ▶ 789 redox-flow batteries
- ▶ Multi-agent system
 - ▶ heat pumps, PVs, CHPs, batteries controlled
 - ▶ demand of households, wind power not controlled
 - ▶ 2710 units for coalition formation
- ▶ Exemplary day: January 31st
- ▶ Coalitions form for hourly products day ahead

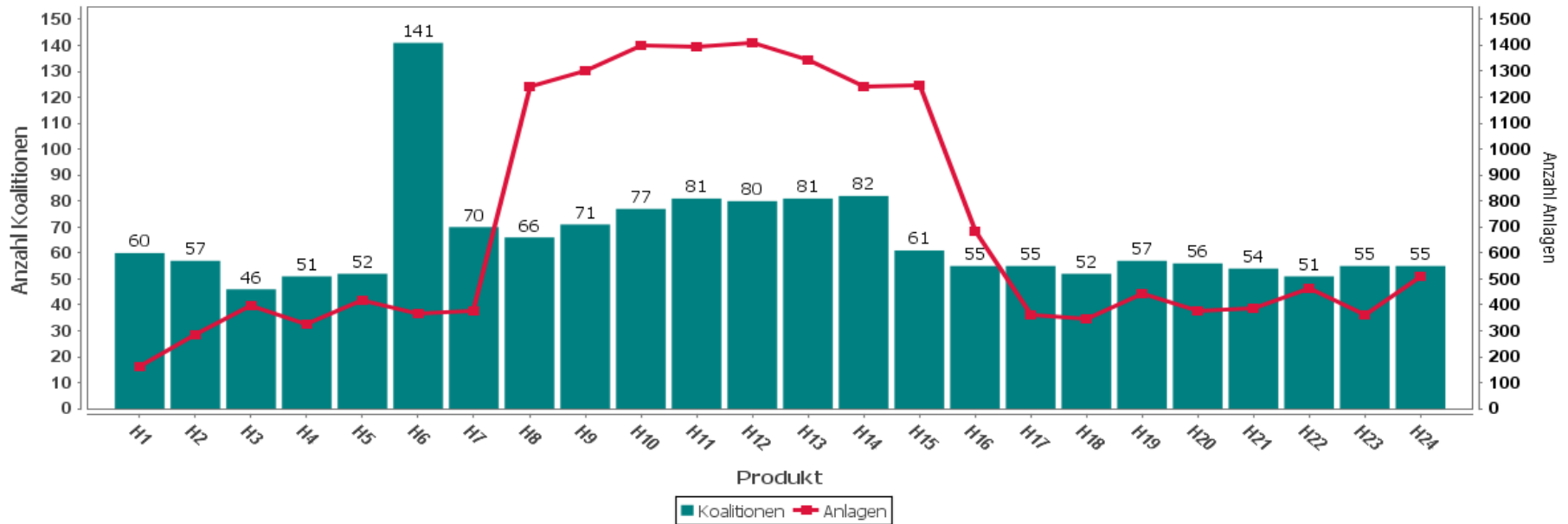
Configuration, Size, and Technology Profile of Coalitions



Configuration, Size, and Technology Profile of Coalitions



Configuration, Size, and Technology Profile of Coalitions



Planning schedules for units in a DVPP

PREDICTIVE SCHEDULING

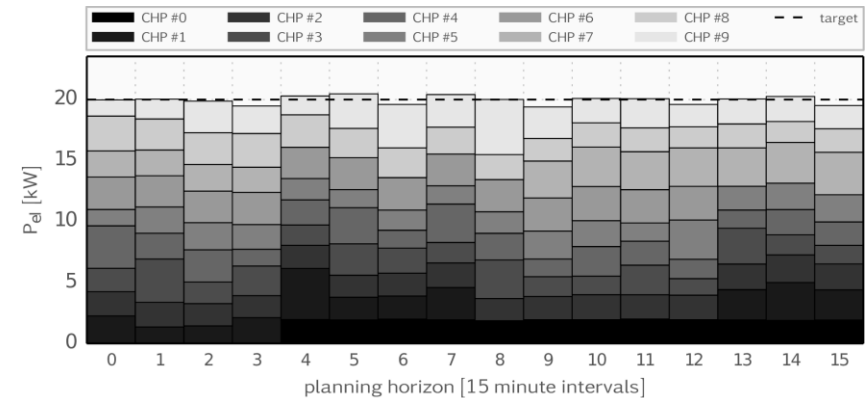
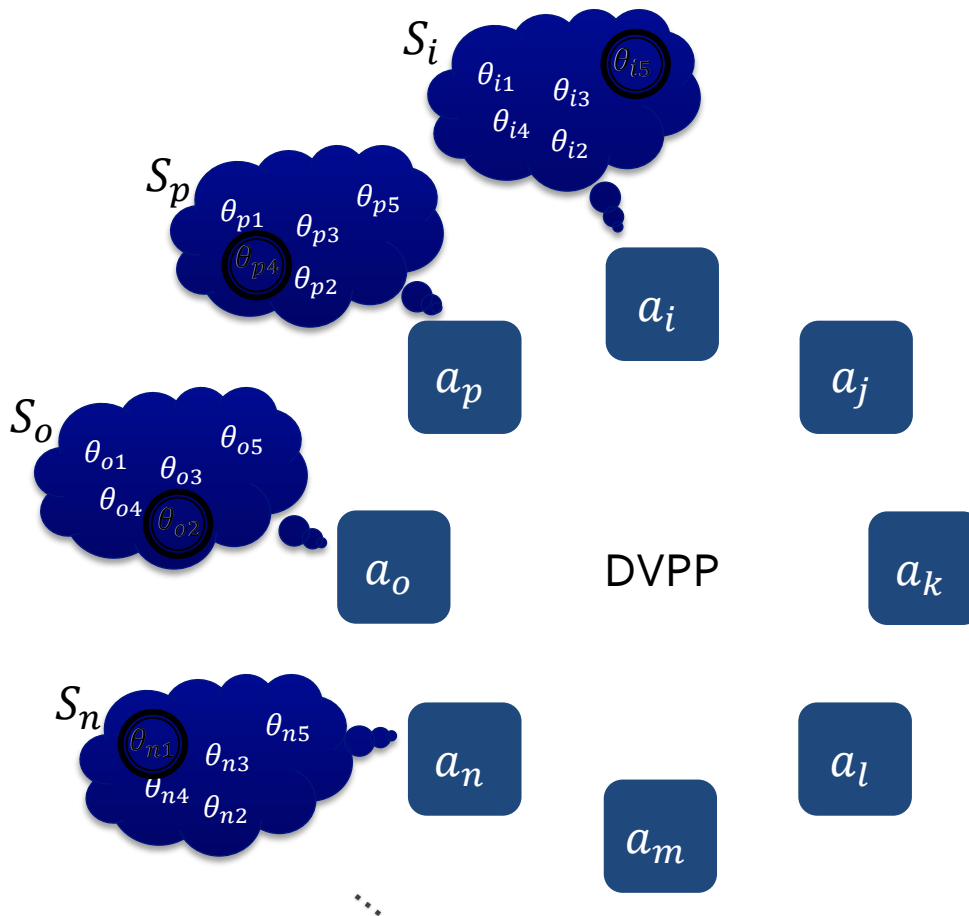
Predictive Planning of Operation Schedules

Goal:

Select an optimal schedule for each unit in a DVPP

- with respect to the agreed power product
- with respect to each unit's private preferences
- in a completely decentralized manner

Predictive Planning as Combinatorial Optimization Problem

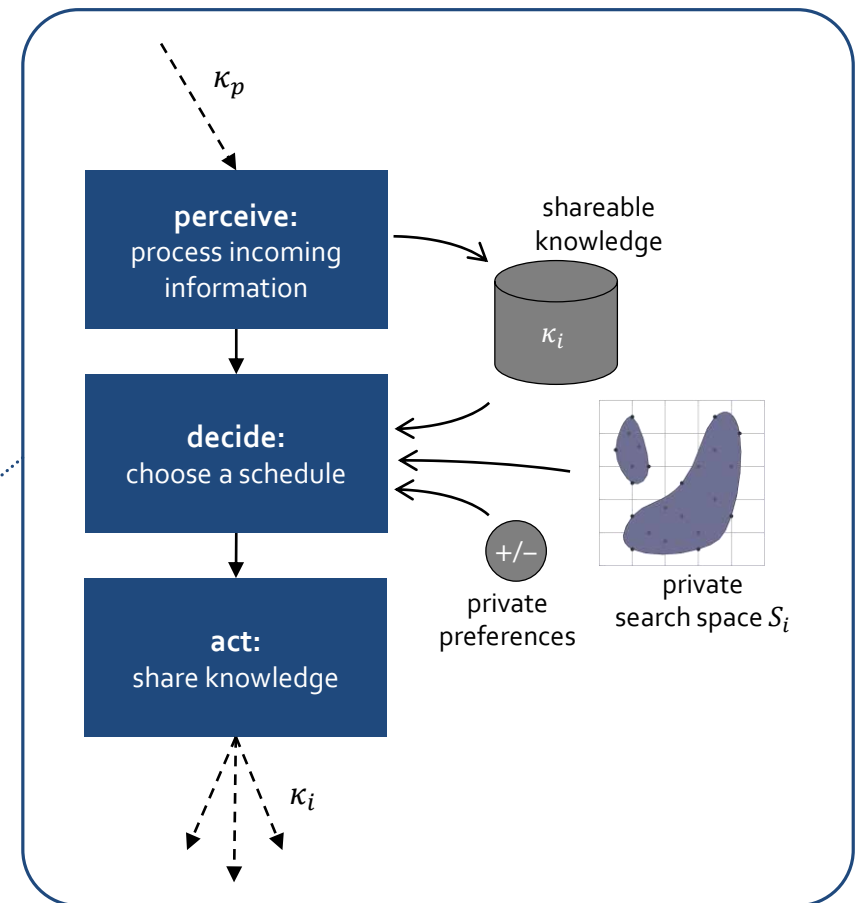
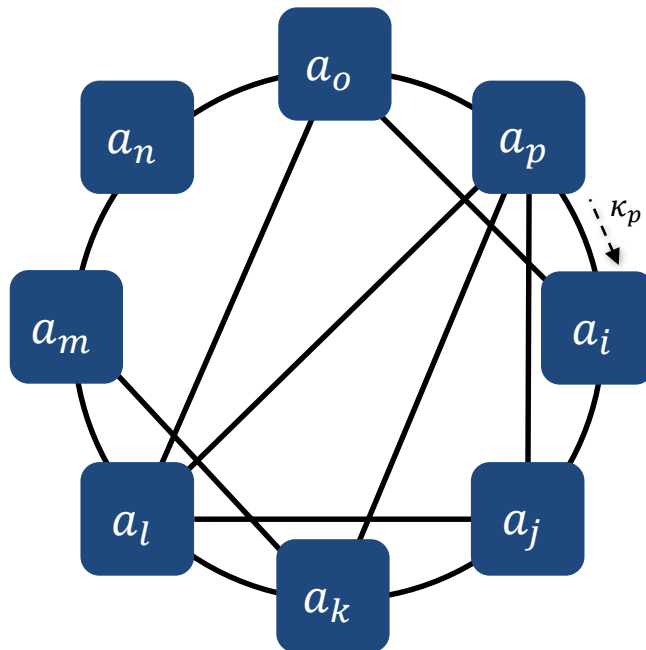


$$\min \left\| \zeta - \sum_{i=1}^{|A|} \sum_{j=1}^{|S_i|} (\theta_{ij} \cdot x_{ij}) \right\|_1$$

subject to $\sum_{j=1}^{|S_i|} x_{ij} = 1, \quad x_{ij} \in \{0, 1\}, \quad i = 1 \dots |A|$

Combinatorial Optimization Heuristic for Distributed Agents

1. Construction of Overlay Communication Topology
2. Agents react asynchronously to exchanged messages



Convergence and termination formally proved using convergence stairs:

- ▶ $\mathcal{A}_1$: Working memories of all agents are complete.
- ▶ $\mathcal{A}_2$: A final solution candidate is found.
- ▶ $\mathcal{A}_3$: Working memories of all agents are identical; heuristic terminates.

After initial setup ($\mathcal{A}_1$), COHDA shows the **anytime** property:

- ▶ Newly produced solutions are always better than older solutions.

Run-Time properties:

- ▶ Best-Case: $\mathcal{O}(3)$
- ▶ Worst-Case: $\mathcal{O}(n^m)$
- ▶ Empirically: $\mathcal{O}(\log m)$

(m = number of agents,
 n = avg. size of search spaces)

Algorithm 1 Behavior of an agent a_i in the COHDA heuristic.

Local attributes:

S_i ▷ Set of feasible schedules
 λ_i ▷ Schedule selection counting variable
 $\kappa_i = (\zeta, \Omega_i, \gamma_i)$ ▷ Working memory: target, system config, solution candidate

Objective functions:

$f(\gamma), f(\Omega)$ ▷ Rates γ or Ω according to the global objective function
 $\text{ACCEPT}(\theta)$ ▷ TRUE, iff θ acceptable according to local objectives

Behavior:

```

1: if receiving message  $\kappa_j$  then
2:    $\zeta_j \leftarrow \zeta \in \kappa_j$  ▷ Extract target profile from message
3:    $\Omega_j \leftarrow \Omega_j \in \kappa_j$  ▷ Extract system configuration from message
4:    $\gamma_j \leftarrow \gamma_j \in \kappa_j$  ▷ Extract solution candidate from message
5:    $\kappa'_j \leftarrow (\zeta, \Omega_j, \gamma_j)$  ▷ Create backup of own working memory
6:    $\omega_i \leftarrow \omega_i \in \Omega_i$  ▷ Extract own schedule selection from  $\Omega_i$ 
7:    $K_i \leftarrow \{a_k \mid \omega_k \in \Omega_i\}$  ▷ Extract agent identifiers from  $\gamma_i$ 
8:    $K_j \leftarrow \{a_k \mid \omega_k \in \Omega_j\}$  ▷ Extract agent identifiers from  $\gamma_j$ 

9:   if  $\zeta = \text{NONE}$  then ▷ Store target profile
10:     $\zeta \leftarrow \zeta_j$ 
11:   end if

12:   for  $\omega_k \in \Omega_i, \omega'_k \in \Omega_j$  do ▷ Update  $\Omega_i$ 
13:     if  $\forall \omega_x \in \Omega_i: x \neq k$  or  $\lambda_k \in \omega'_k > \lambda_k \in \omega_k$  then
14:        $\Omega_i \leftarrow (\Omega_i \setminus \{\omega_k\}) \cup \{\omega'_k\}$ 
15:     end if
16:   end for

17:   if  $K_i \subset K_j$  then ▷ Replace/Extend solution candidate...
18:      $\gamma_i \leftarrow \gamma_j$  ▷ ...if the new one is larger
19:   else if  $K_i \not\subset K_j$  then ▷ ...or if it contains entries from unknown agents
20:      $\gamma_i \leftarrow (a_i, \{\omega_k \in \gamma_i\} \cup \{\omega_k \mid a_k \in K_j - K_i\})$ 
21:   else if  $f(\gamma_j) > f(\gamma_i)$  then ▷ ...or if it rates better
22:      $\gamma_i \leftarrow \gamma_j$ 
23:   else if  $f(\gamma_j) = f(\gamma_i)$  and  $a_x \in \gamma_j > a_y \in \gamma_i$  then ▷ ...or to break ties
24:      $\gamma_i \leftarrow \gamma_j$ 
25:   end if

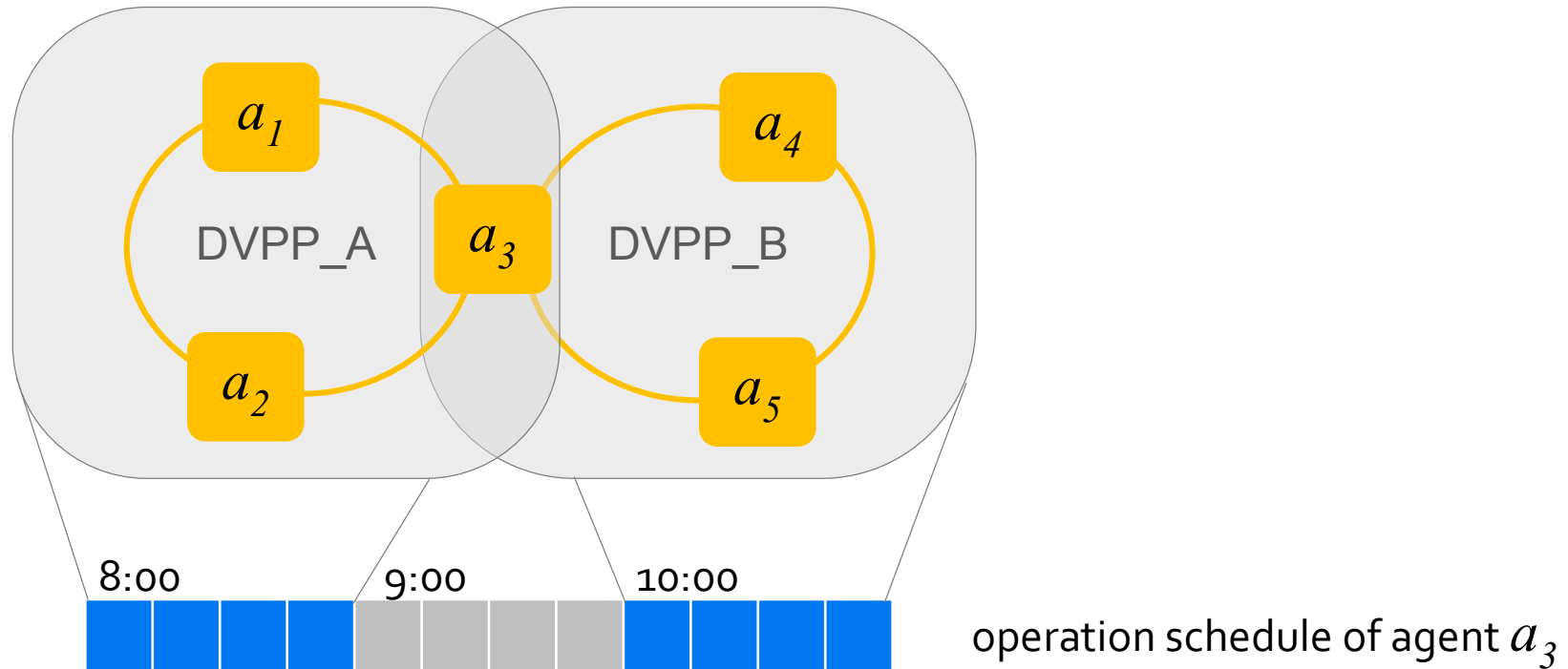
26:    $\omega'_i \leftarrow \omega_i \in \gamma_i$  ▷ Extract own schedule selection from  $\gamma_i$ 
27:   if  $\exists \theta \in S_i: f(\{(\Omega_i \setminus \{\omega_i\}) \cup \{(a_i, \theta, \lambda_i + 1)\}\}) > f(\gamma_i)$  and  $\text{ACCEPT}(\theta)$  then ▷ Better schedule found?
28:      $\omega'_i \leftarrow (a_i, \theta, \lambda_i + 1)$  ▷ Create new schedule selection
29:      $\gamma_i \leftarrow (a_i, (\Omega_i \setminus \{\omega_i\}) \cup \{\omega'_i\})$  ▷ Create new solution candidate
30:   else if  $\omega_i \neq \omega'_i$  then ▷ Historical schedule chosen?
31:      $\omega'_i \leftarrow (a_i, \theta \in \omega'_i, \lambda_i + 1)$  ▷ Create new schedule selection using the historical schedule
32:   end if
33:    $\Omega_i \leftarrow (\Omega_i \setminus \{\omega_i\}) \cup \{\omega'_i\}$  ▷ Update  $\Omega_i$  with new schedule selection

34:   if  $(\zeta, \Omega_i, \gamma_i) \neq \kappa'_i$  then ▷ Publish working memory to neighbours
35:     send  $(\zeta, \Omega_i, \gamma_i)$  to neighbours
36:   end if

37: end if

```

COHDA for multiple DVPPs

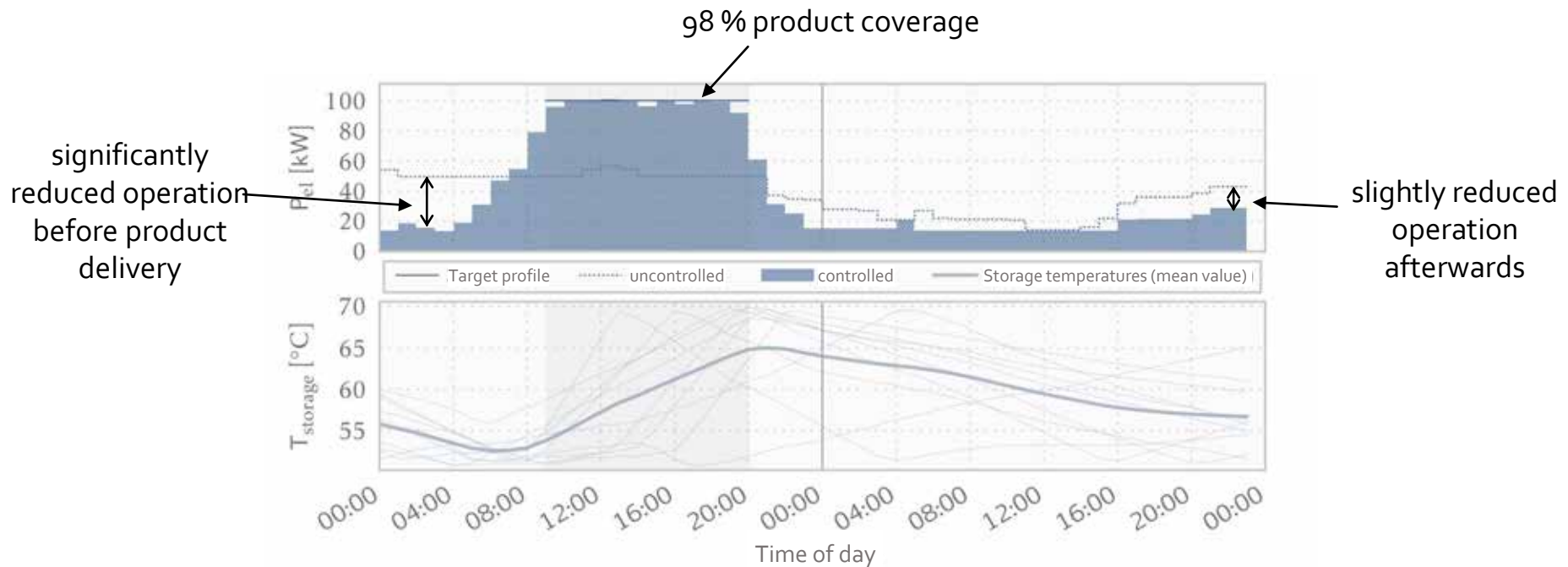


Evaluation of solution candidate only considering
DVPP under negotiation

Example: EPEX Spot

Scenario:

- ▶ 10 CHP (Vaillant EcoPower 20.0)
- ▶ EPEX Spot, Block Product "Peak Load" (100 kW, 9 am – 8 pm)

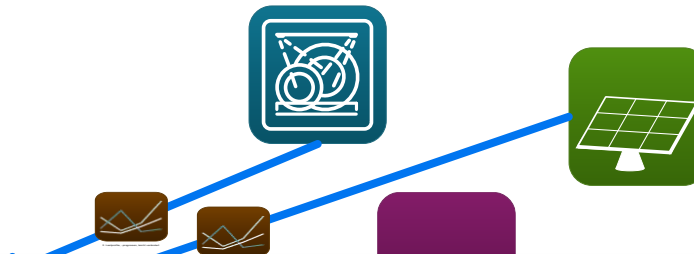


Coping large sets of feasible schedules

REPRESENTING FLEXIBILITIES

Goals of representing flexibilities

Essential for the acceptance of planning algorithms: Feasibility of solutions!



How can individual flexibilities, secondary goals, preferences, cost, etc. be considered in optimization?



Representation of flexibilities

Goal:

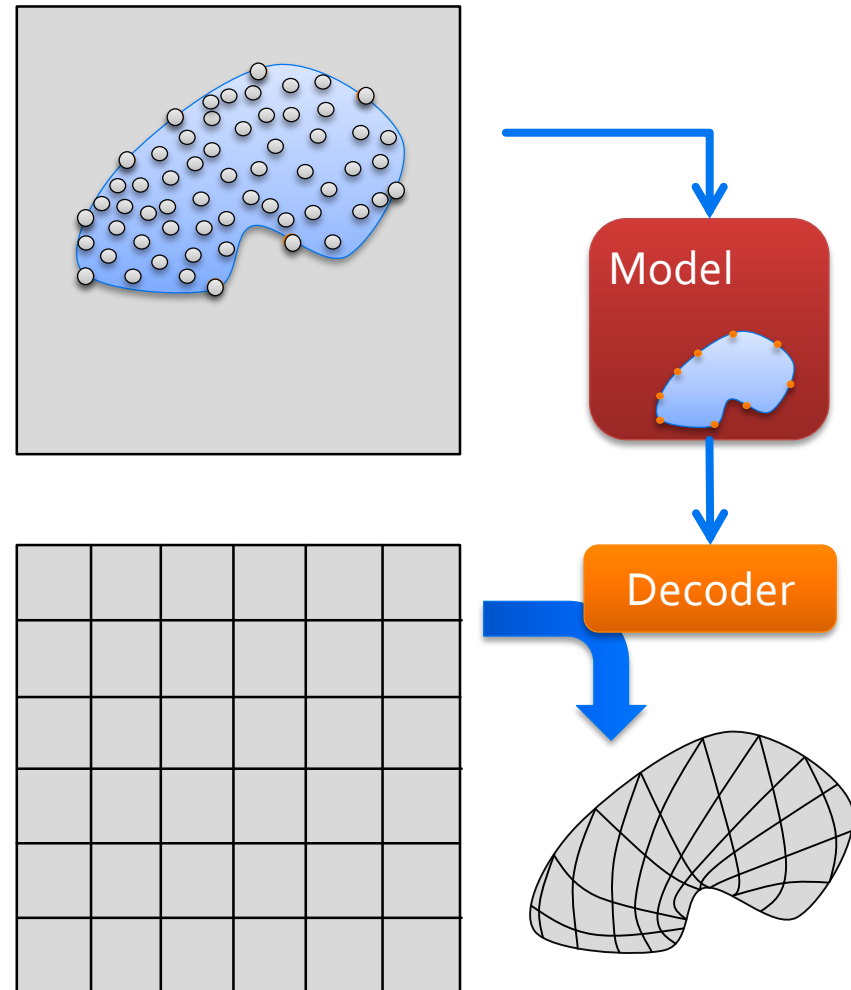
Representation of a unit's search space must be

- independent of individual configuration of units
- efficiently accessible for optimisation processes
- allow the representation of additional features
- hide private information

Representation of schedules / flexibilities

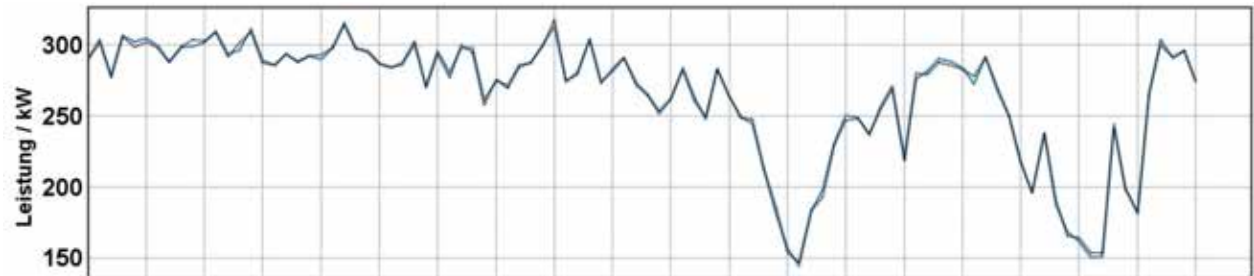
Process chain

- ▶ Samples of feasible schedules simulated as patterns for search space
- ▶ Modelling of search space by SVM classifier
- ▶ Derivation of decoder function to map arbitrary schedules into the space of feasible schedules
 - ▶ Unconstrained optimisation

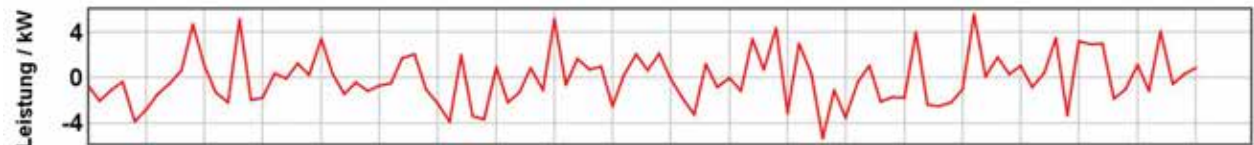


Exemplary result

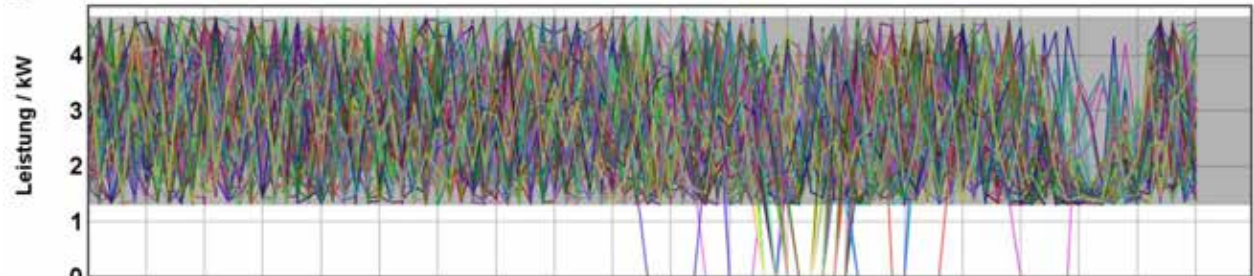
Goal function and
result of optimization



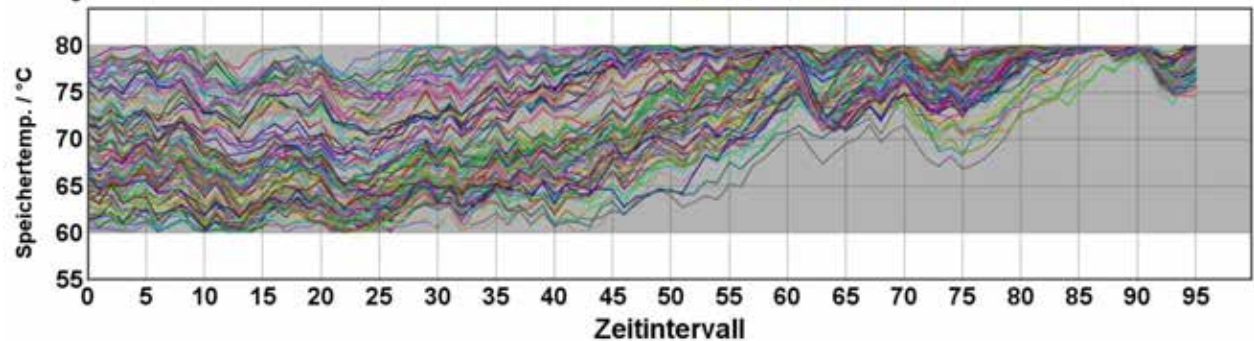
Residual error



Feasible schedules for
100 CHPs



Feasible temperature
levels for the heat
buffers of the CHPs



Reacting to events during product delivery

CONTINUOUS SCHEDULING

Continuous Energy Scheduling

Market

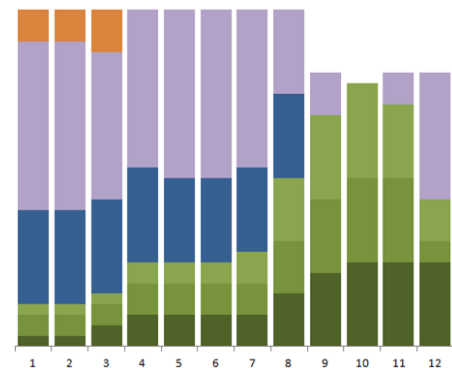
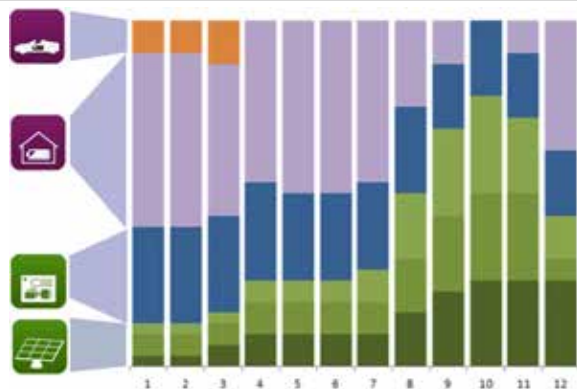
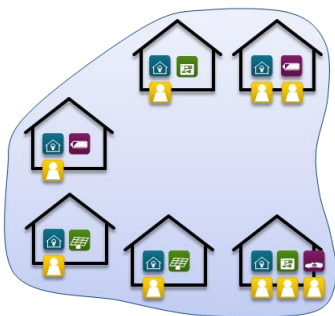


Product

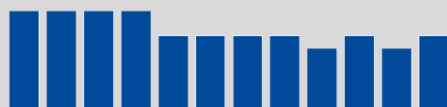


?

Coalition



Unit



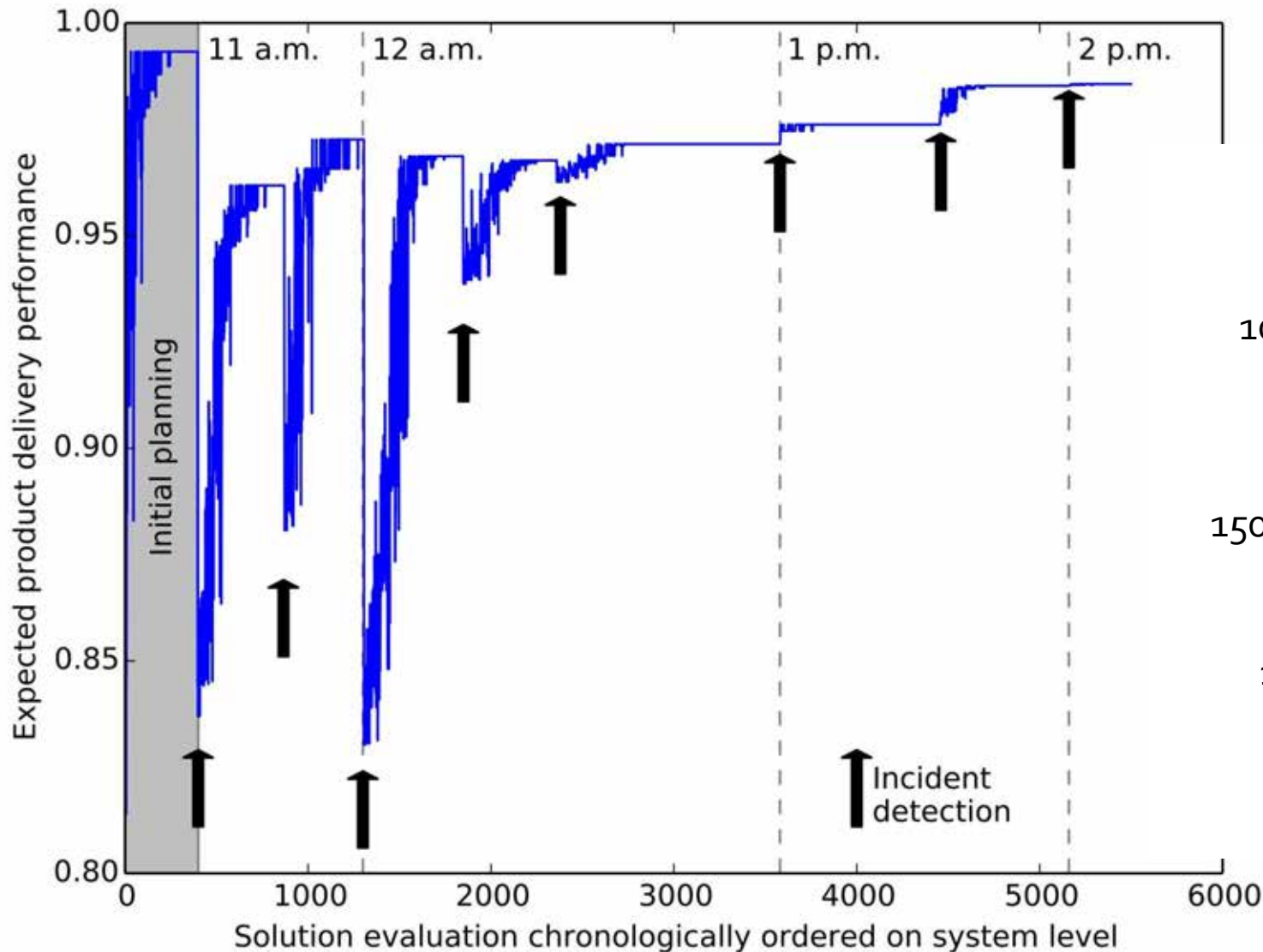
CHP schedule



Incident types

- Breakdown
- Prognosis fault
- System service

Results: Rescheduling



DVPP:

20 CHP(4,7 kW_{el})
20 PV (10 kW_p)
10 heat pumps(5kW)

Product:

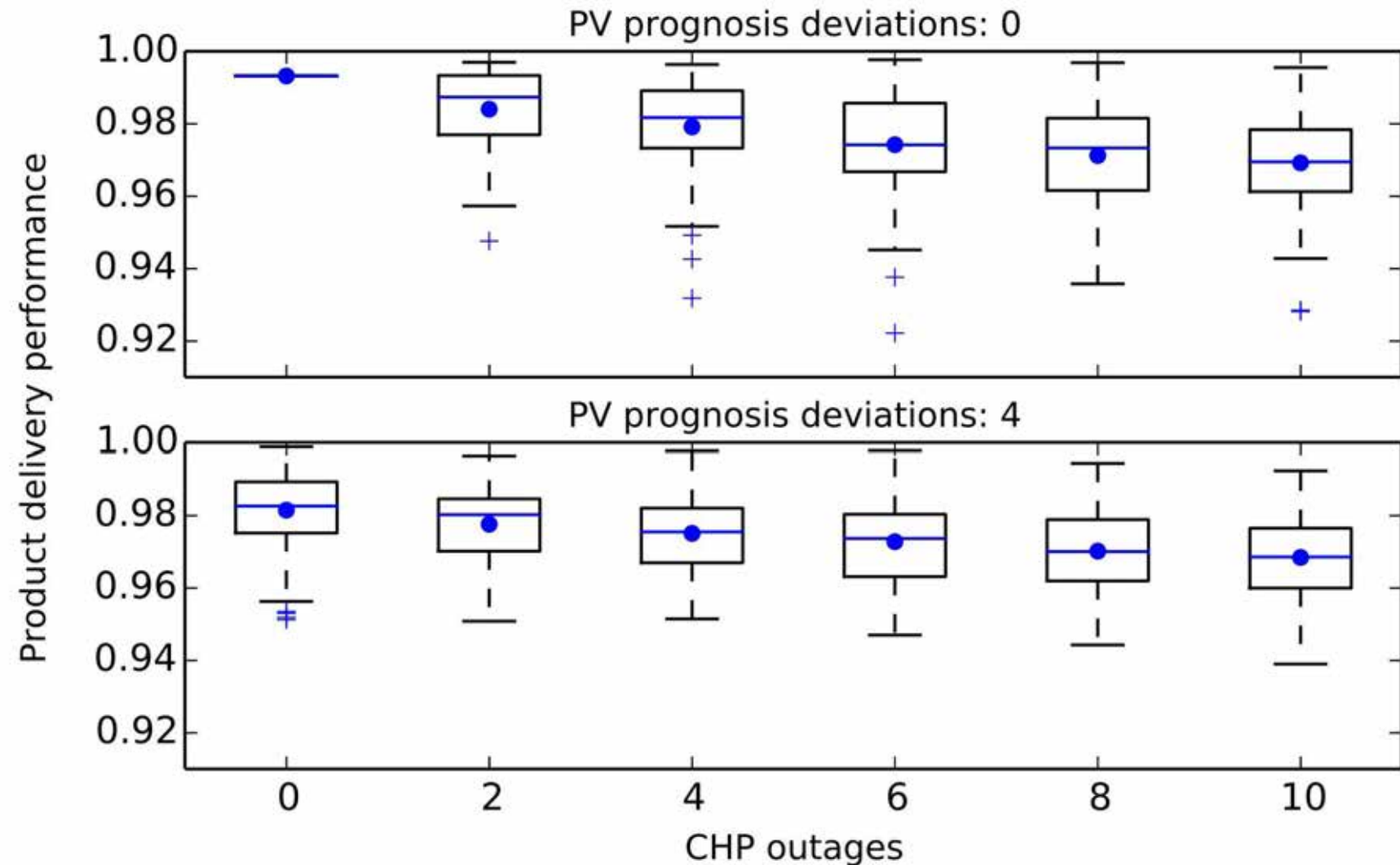
4 h
150 kWh for each hour

Incidents:

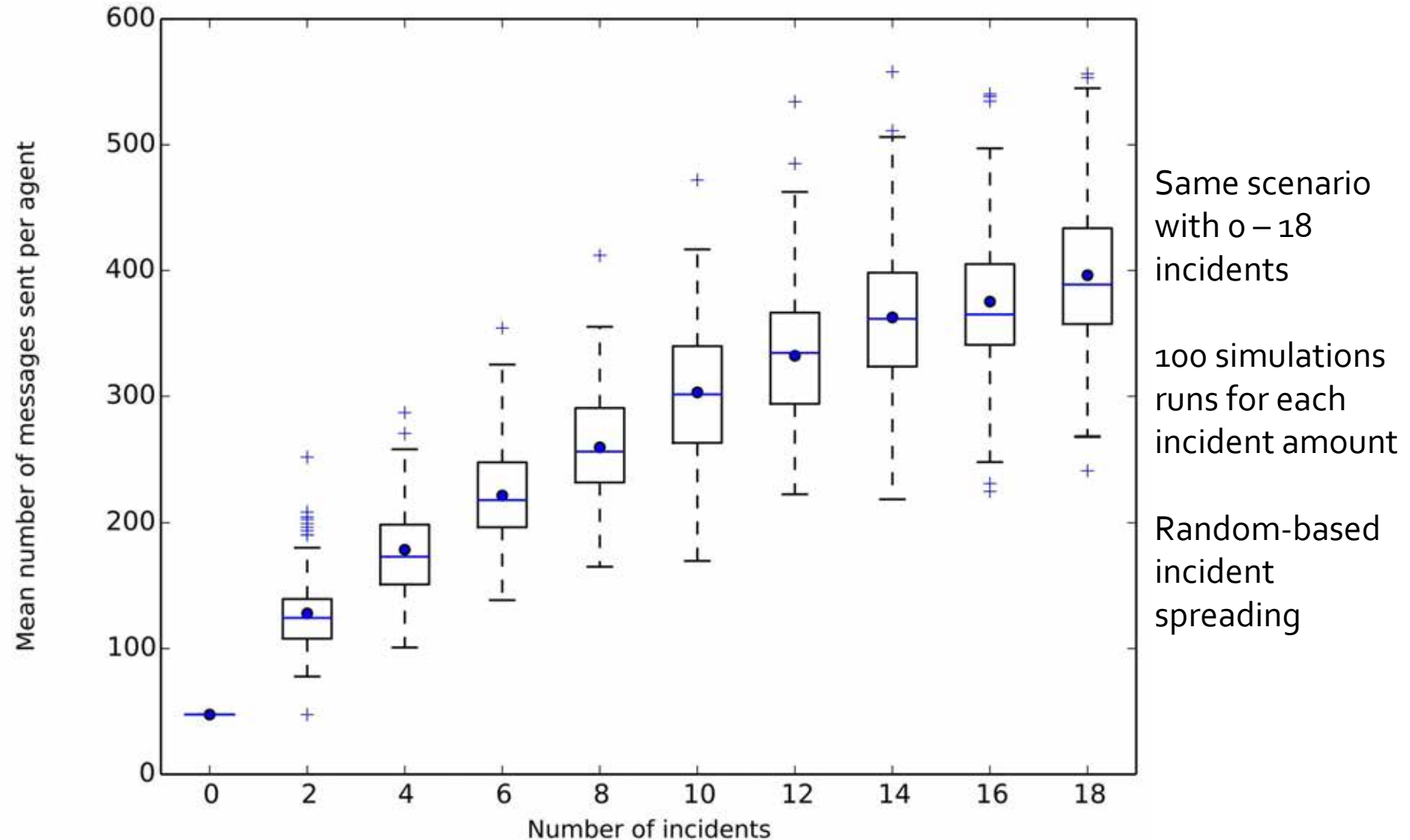
10 CHP shut-downs
(30min)

Day in spring

Results: Rescheduling



Message overhead



CONCLUSION

Completely self-organized virtual power plant

- ▶ coalition formation
- ▶ predictive scheduling
- ▶ reactive scheduling
- ▶ abstract representation of flexibilities

Smart grid algorithm engineering

Methodical challenges

- ▶ Engineering approach for self-organization methods
- ▶ 'Well-behaviour' of distributed self-organizing methods
 - controlled self-organization as in organic computing
- ▶ Security aspects in the engineering process
- ▶ Representative evaluation scenarios
- ▶ Computational effort for sensitivity analysis
- ▶ ...



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More information required?

<http://smarnord.de/downloads/SmartNordFinalReport.pdf>